

Advanced Functional Programming, 1DL450

Lecture 10, 2012-12-10

Cons T Åhs

Monads revisited

- ▶ Monads combine pure functional computation with side effects.
- ▶ Sequencing and order of computation are important for the side effects.
- ▶ With monads it is much more clear which parts contain side effects/are dependent on order and which are purely functional.
- ▶ For I/O, graphics etc it is “obvious” that we need sequential computation.
- ▶ It is less obvious how sequentiality is described formally, especially when combined with a lazy language.
- ▶ Using the formal description one can use monads to describe more types of computation than just sequential.
- ▶ A monad contains (at least) two operators
 - ▶ *bind* or $>>=$ used for sequential operation and (optional) binding of name
 - ▶ *return* for extracting values from the monad

Monad Axioms

- ▶ *Return* acts a “neutral” (or identity; $x + 0 = 0 + x = x$) element of *bind*

$$\begin{aligned}(\mathbf{return\ x}) \gg= f &\equiv f\ x \\ m \gg= \mathbf{return} &\equiv m\end{aligned}$$

- ▶ Binding two functions in succession is the same as binding one function that can be determined from them
 - ▶ This is a composition property

$$(m \gg= f) \gg= g \equiv m \gg= (\backslash x \rightarrow (f\ x \gg= g))$$

- ▶ Optional *zero* value ($x * 0 = 0 * x = 0$)

$$mzero \gg= f \equiv mzero$$

$$m \gg= (\backslash x \rightarrow mzero) \equiv mzero$$

Maybe Monad

- ▶ Computations might fail (in general)
- ▶ Can we express and propagate failed computations in an easy way?
 - ▶ Compare to exceptions; skip rest of computation if one fails

```
data Maybe a = The a | Nil deriving Show
```

```
instance Monad Maybe where
  (The x) >>= f = f x
  Nil >>= _ = Nil
  return x = The x
```

```
madd x y =
  do
    x1 <- x
    y1 <- y
    return (x1 + y1)
```

```
*Main> :type madd
madd :: (Monad m, Num b) => m b -> m b -> m b
*Main> madd (The 2) (The 3)
The 5
*Main> madd (The 2) Nil
Nil
```

Maybe Monad

- ▶ Be aware the sequencing/bind “unpacks” values inside monads
- ▶ The following two definitions of monadic add are not equivalent
 - ▶ observe the very different types!

```
madd x y =  
  do  
    x1 <- x  
    y1 <- y  
    return (x1 + y1)
```

```
madd2 x y =  
  do return (x + y)
```

```
*Main> :type madd  
madd :: (Monad m, Num b) => m b -> m b -> m b  
*Main> :type madd2  
madd2 :: (Monad m, Num a) => a -> a -> m a
```

List Monad

- ▶ Use a monad for a set of possible results
- ▶ Note that it makes sense to read “<-” as “belongs to” or “in”
- ▶ Very much like list comprehension

```
instance Monad [] where
  -- single possible result
  return a = [a]
  -- apply to each element and flatten results
  xs >>= f = concat (map f xs)
  -- failure is no results
  fail s = []
```

```
Prelude> do x<-[1,2,3]; y<-[4,5]; return (x*y)
[4,5,8,10,12,15]
Prelude> do x<-[1,2,3]; y<-[4,5]; if x > 1 then return (x*y)
else fail ""
[8,10,12,15]
```

Other Monads

- ▶ State monad
 - ▶ our previous cell example was a state monad
- ▶ Environment Monad
 - ▶ compute with values from a (non local) context
- ▶ Writer monad
 - ▶ perform computations while writing “output” on the side
- ▶ See Wikipedia page for more details.