

Code Generation

The Main Idea of Today's Lecture

We can emit stack-machine-style code for expressions via recursion

(We will use MIPS assembly as our target language)

Lecture Outline

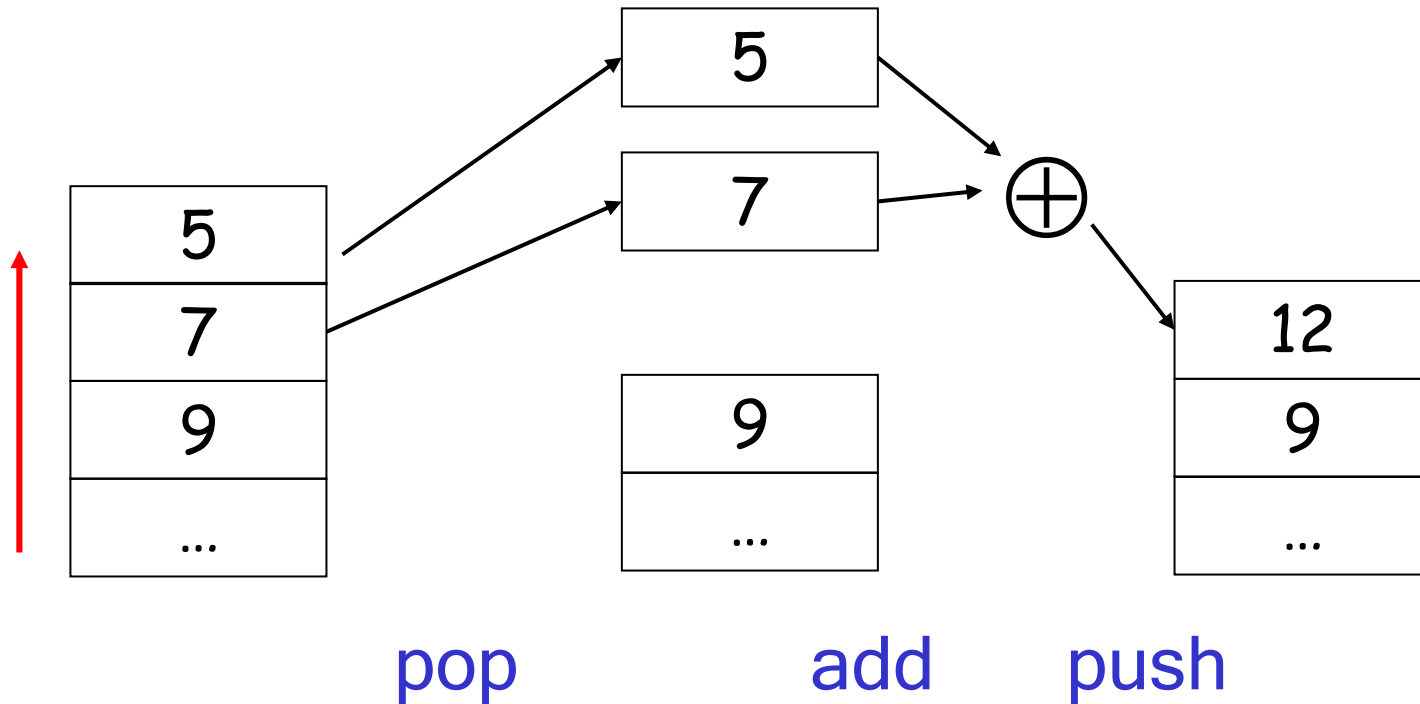
- What are stack machines?
- The MIPS assembly language
- A simple source language (“Mini Bar”)
- A stack machine implementation of the simple language

Stack Machines

- A simple evaluation model
- No variables or registers
- A stack of values for intermediate results
- Each **instruction**:
 - Takes its operands from the top of the stack
 - Removes those operands from the stack
 - Computes the required operation on them
 - Pushes the result onto the stack

Example of Stack Machine Operation

The addition operation on a stack machine



Example of a Stack Machine Program

- Consider two instructions
 - `push i` - place the integer `i` on top of the stack
 - `add` - pop topmost two elements, add them and put the result back onto the stack
- A program to compute $7 + 5$:
 - `push 7`
 - `push 5`
 - `add`

Why Use a Stack Machine?

- Each operation takes operands from the same place and puts results in the same place
- This means a uniform compilation scheme
- And therefore a simpler compiler

Why Use a Stack Machine?

- Location of the operands is implicit
 - Always on the top of the stack
- No need to specify operands explicitly
- No need to specify the location of the result
- Instruction is "add" as opposed to "add r_1, r_2 " (or "add $r_d r_{i1} r_{i2}$ ")
 - ⇒ Smaller encoding of instructions
 - ⇒ More compact programs
- This is one of the reasons why Java Bytecode uses a stack evaluation model

Optimizing the Stack Machine

- The `add` instruction does 3 memory operations
 - Two reads and one write to the stack
 - The top of the stack is frequently accessed
- Idea: keep the top of the stack in a dedicated register (called the "accumulator")
 - Register accesses are faster (why?)
- The "`add`" instruction is now
$$\text{acc} \leftarrow \text{acc} + \text{top_of_stack}$$
 - Only one memory operation!

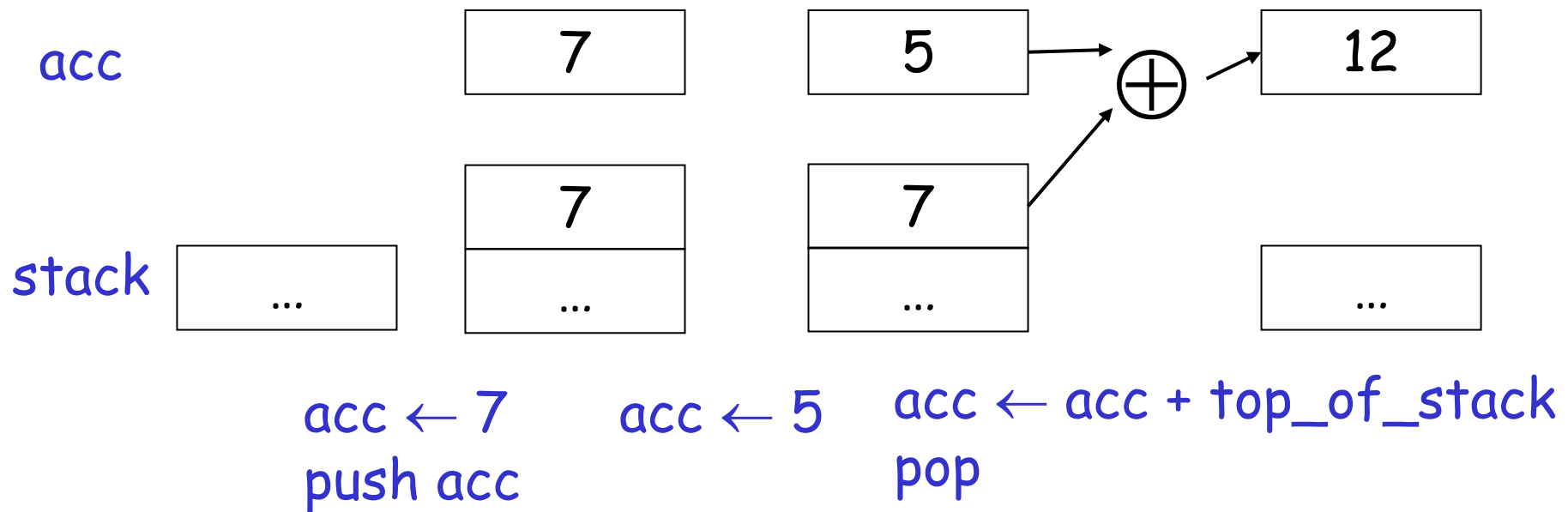
Stack Machine with Accumulator

Invariants

- The result of computing an expression is always placed in the accumulator
- For an operation $op(e_1, \dots, e_n)$ compute each e_i and then push the accumulator (= the result of evaluating e_i) onto the stack
- After the operation pop $n-1$ values
- After computing an expression the stack is as before

Stack Machine with Accumulator: Example

Compute $7 + 5$ using an accumulator



A Bigger Example: $3 + (7 + 5)$

Code	Acc	Stack
<code>acc ← 3</code>	3	<init>
<code>push acc</code>	3	3, <init>
<code>acc ← 7</code>	7	3, <init>
<code>push acc</code>	7	7, 3, <init>
<code>acc ← 5</code>	5	7, 3, <init>
<code>acc ← acc + top_of_stack</code>	12	7, 3, <init>
<code>pop</code>	12	3, <init>
<code>acc ← acc + top_of_stack</code>	15	3, <init>
<code>pop</code>	15	<init>

Notes

- It is very important that the stack is preserved across the evaluation of a subexpression
 - Stack before the evaluation of $7 + 5$ is $3, \langle \text{init} \rangle$
 - Stack after the evaluation of $7 + 5$ is $3, \langle \text{init} \rangle$
 - The first operand is on top of the stack

From Stack Machines to MIPS

- The compiler generates code for a stack machine with accumulator
- We want to run the resulting code on the MIPS processor (or simulator)
- We simulate the stack machine instructions using MIPS instructions and registers

Simulating a Stack Machine on the MIPS...

- The accumulator is kept in MIPS register $\$a0$
- The stack is kept in memory
- The stack grows towards lower addresses
 - Standard convention on the MIPS architecture
- The address of the next location on the stack is kept in MIPS register $\$sp$
 - Guess: what does "sp" stand for?
 - The top of the stack is at address $\$sp + 4$

MIPS Assembly

MIPS architecture

- Prototypical Reduced Instruction Set Computer (RISC) architecture
- Arithmetic operations use registers for operands and results
- Must use `load` and `store` instructions to use operands and store results in memory
- 32 general purpose registers (32 bits each)
 - We will use `$sp`, `$a0` and `$t1` (a temporary register)

Read the SPIM documentation for more details

A Sample of MIPS Instructions

- lw reg_1 offset(reg_2) "load word"
 - Load 32-bit word from address $reg_2 + \text{offset}$ into reg_1
- add reg_1 reg_2 reg_3
 - $reg_1 \leftarrow reg_2 + reg_3$
- sw reg_1 offset(reg_2) "store word"
 - Store 32-bit word in reg_1 at address $reg_2 + \text{offset}$
- addiu reg_1 reg_2 imm "add immediate"
 - $reg_1 \leftarrow reg_2 + \text{imm}$
 - "u" means overflow is not checked
- li reg imm "load immediate"
 - $reg \leftarrow \text{imm}$

MIPS Assembly: Example

- The stack-machine code for $7 + 5$ in MIPS:

$acc \leftarrow 7$
push acc

$acc \leftarrow 5$
 $acc \leftarrow acc + \text{top_of_stack}$

pop

```
li $a0 7
sw $a0 0($sp)
addiu $sp $sp -4
li $a0 5
lw $t1 4($sp)
add $a0 $a0 $t1
addiu $sp $sp 4
```

- We now generalize this to a simple language...

A Small Language

- A language with only integers and integer operations (“Mini Bar”)

$$P \rightarrow F P \mid F$$
$$F \rightarrow \text{id}(\text{ARGS}) \text{ begin } E \text{ end}$$
$$\text{ARGS} \rightarrow \text{id}, \text{ARGS} \mid \text{id}$$
$$E \rightarrow \text{int} \mid \text{id} \mid \text{if } E_1 = E_2 \text{ then } E_3 \text{ else } E_4 \\ \mid E_1 + E_2 \mid E_1 - E_2 \mid \text{id}(\text{ES})$$
$$\text{ES} \rightarrow E, \text{ES} \mid E$$

A Small Language (Cont.)

- The first function definition f is the "main" routine
- Running the program on input i means computing $f(i)$
- Program for computing the Fibonacci numbers:

```
fib(x)
```

```
begin
```

```
    if x = 1 then 0 else
```

```
    if x = 2 then 1 else fib(x - 1) + fib(x - 2)
```

```
end
```

Code Generation Strategy

- For each expression e we generate MIPS code that:
 - Computes the value of e in $\$a0$
 - Preserves $\$sp$ and the contents of the stack
- We define a code generation function $cgen(e)$ whose result is the code generated for e
 - $cgen(e)$ will be recursive

Code Generation for Constants

- The code to evaluate an integer constant simply copies it into the accumulator:

`cgen(int) = li $a0 int`

- Note that this also preserves the stack, as required

Code Generation for Addition

```
cgen( $e_1 + e_2$ ) =  
    cgen( $e_1$ )           ; $a0 ← value of  $e_1$   
    sw $a0 0($sp)        ; push that value  
    addiu $sp $sp -4     ; onto the stack  
    cgen( $e_2$ )           ; $a0 ← value of  $e_2$   
    lw $t1 4($sp)       ; grab value of  $e_1$   
    add $a0 $t1 $a0      ; do the addition  
    addiu $sp $sp 4     ; pop the stack
```

Possible optimization:

Put the result of e_1 directly in register \$t1?

Code Generation for Addition: Wrong Attempt!

Optimization: Put the result of e_1 directly in $\$t1$?

```
cgen( $e_1 + e_2$ ) =  
    cgen( $e_1$ )           ;  $\$a0 \leftarrow$  value of  $e_1$   
    move  $\$t1$   $\$a0$        ; save that value in  $\$t1$   
    cgen( $e_2$ )           ;  $\$a0 \leftarrow$  value of  $e_2$   
                               ; may clobber  $\$t1$   
    add  $\$a0$   $\$t1$   $\$a0$    ; perform the addition
```

Try to generate code for : $3 + (7 + 5)$

Code Generation Notes

- The code for $e_1 + e_2$ is a template with “holes” for code for evaluating e_1 and e_2
- Stack machine code generation is recursive
- Code for $e_1 + e_2$ consists of code for e_1 and e_2 glued together
- Code generation can be written as a recursive-descent of the AST
 - At least for (arithmetic) expressions

Code Generation for Subtraction and Constants

New instruction: `sub reg1 reg2 reg3`

Implements $\text{reg}_1 \leftarrow \text{reg}_2 - \text{reg}_3$

`cgen(e1 - e2) =`

<code>cgen(e₁)</code>	<code>; \$a0 ← value of e₁</code>
<code>sw \$a0 0(\$sp)</code>	<code>; push that value</code>
<code>addiu \$sp \$sp -4</code>	<code>; onto the stack</code>
<code>cgen(e₂)</code>	<code>; \$a0 ← value of e₂</code>
<code>lw \$t1 4(\$sp)</code>	<code>; grab value of e₁</code>
<code>sub \$a0 \$t1 \$a0</code>	<code>; do the subtraction</code>
<code>addiu \$sp \$sp 4</code>	<code>; pop the stack</code>

Code Generation for Conditional

- We need flow control instructions
- New MIPS instruction: `beq reg1 reg2 label`
 - Branch to `label` if `reg1 = reg2`
- New MIPS instruction: `j label`
 - Unconditional jump to `label`

Code Generation for If (Cont.)

cgen(if $e_1 = e_2$ then e_3 else e_4) =

cgen(e_1)

sw \$a0 0(\$sp)

addiu \$sp \$sp -4

cgen(e_2)

lw \$t1 4(\$sp)

addiu \$sp \$sp 4

beq \$a0 \$t1 true_branch

false_branch:

cgen(e_4)

j end_if

true_branch:

cgen(e_3)

end_if:

Meet The Activation Record

- Code for function calls and function definitions depends on the layout of the activation record (or "AR")
- A very simple AR suffices for this language:
 - The result is always in the accumulator
 - No need to store the result in the AR
 - The activation record holds actual parameters
 - For $f(x_1, \dots, x_n)$ push the arguments x_n, \dots, x_1 onto the stack
 - These are the only variables in this language

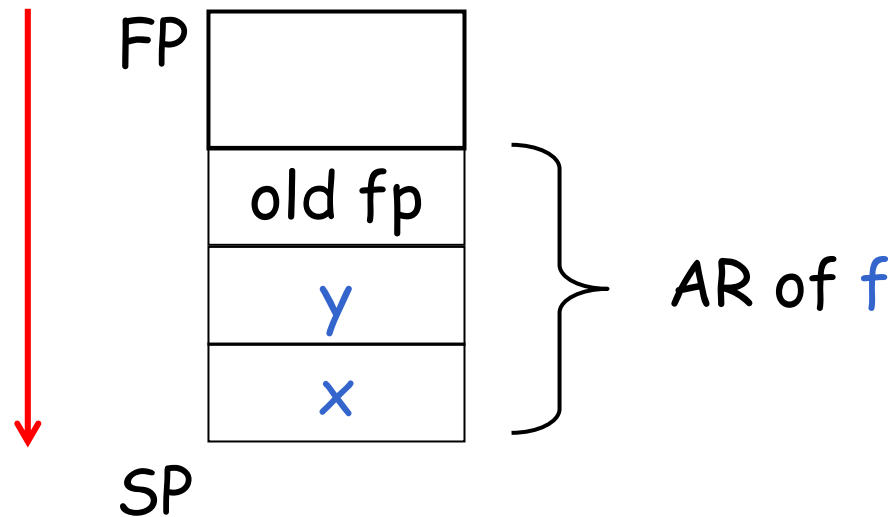
Meet The Activation Record (Cont.)

- The stack discipline guarantees that on function exit, `$sp` is the same as it was before the args got pushed (i.e., before function call)
- We need the return address
- It's also handy to have a pointer to the current activation
 - This pointer lives in register `$fp` (frame pointer)
 - Reason for frame pointer will be clear shortly (at least I hope!)

Layout of the Activation Record

Summary: For this language, an AR with the caller's frame pointer, the actual parameters, and the return address suffices

Picture: Consider a call to $f(x,y)$, the AR will be:



Code Generation for Function Call

- The calling sequence is the sequence of instructions (of both *caller* and *callee*) to set up a function invocation
- New instruction: `jal label`
 - Jump to `label`, save address of next instruction in special register `$ra`
 - On other architectures the return address is stored on the stack by the "`call`" instruction

Code Generation for Function Call (Cont.)

```
cgen(f(e1,...,en)) =  
  sw $fp 0($sp)  
  addiu $sp $sp -4  
  cgen(en)  
  sw $a0 0($sp)  
  addiu $sp $sp -4  
  ...  
  cgen(e1)  
  sw $a0 0($sp)  
  addiu $sp $sp -4  
  jal f_entry
```

- The caller saves the value of the frame pointer
- Then it pushes the actual parameters in reverse order
- The caller's **jal** puts the return address in register **\$ra**
- The AR so far is $4*n+4$ bytes long

Code Generation for Function Definition

- New MIPS instruction: `jr reg`
 - Jump to address in register `reg`

`cgen(f(x1,...,xn) begin e end) =`

`f_entry:`

`move $fp $sp`

`sw $ra 0($sp)`

`addiu $sp $sp -4`

`cgen(e)`

`lw $ra 4($sp)`

`addiu $sp $sp frame_size`

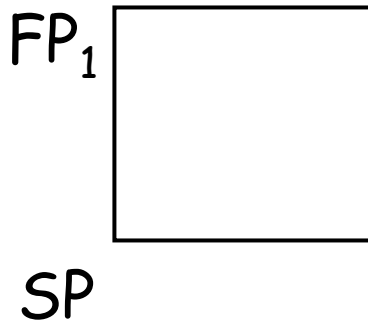
`lw $fp 0($sp)`

`jr $ra`

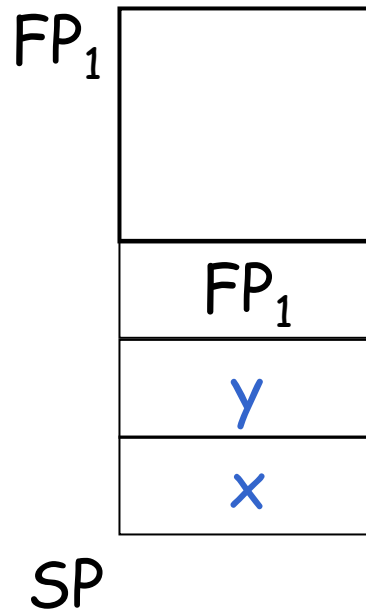
- Note: The frame pointer points to the top, not bottom of the frame
- Callee saves old return address, evaluates its body, pops the return address, pops the arguments, and then restores `$fp`
- `frame_size = 4*n + 8`

Calling Sequence: Example for $f(x,y)$

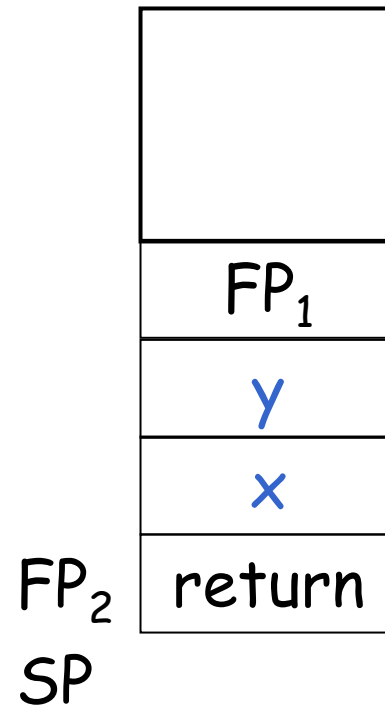
Before call



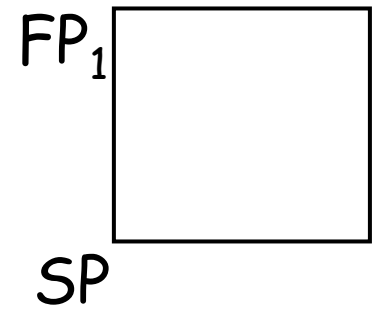
On entry



After body



After call



Code Generation for Variables/Parameters

- Variable references are the last construct
- The “variables” of a function are just its parameters
 - They are all in the AR
 - Pushed by the caller
- Problem: Because the stack grows when intermediate results are saved, the variables are not at a fixed offset from $\$sp$

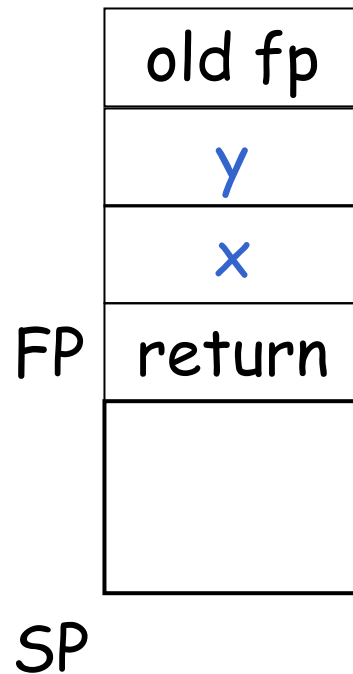
Code Generation for Variables/Parameters

- Solution: use the frame pointer
 - Always points to the return address on the stack
 - Since it does not move, it can be used to find the variables
- Let x_i be the i^{th} ($i = 1, \dots, n$) formal parameter of the function for which code is being generated

$\text{cgen}(x_i) = \text{lw } \$a0, \text{offset}(\$fp)$ ($\text{offset} = 4 * i$)

Code Generation for Variables/Parameters

- Example: For a function $f(x,y)$ `begin e end` the activation and frame pointer are set up as follows (when evaluating e):



- x is at $fp + 4$
- y is at $fp + 8$

Activation Record & Code Generation Summary

- The activation record must be designed together with the code generator
- Code generation can be done by recursive traversal of the AST

Discussion

- Production compilers do different things
 - Emphasis is on keeping values (esp. current stack frame) in registers
 - Intermediate results are laid out in the AR, not pushed and popped from the stack
 - As a result, code generation is often performed in synergy with register allocation

Next time: code generation for temporaries and a deeper look into parameter passing mechanisms