

Exam for Real Time Systems

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Important Instructions:

1. No course material or computer/calculator are allowed, only a pen and a dictionary.
2. Please mark which course you are registered for:

5hp (1DT063)

You need to solve problems 1–6 only.

10hp (1DT004)

You need to solve *all* problems.

3. For each problem, a number of choices/statements are given, that may be correct or wrong.

You are asked to mark only the **correct** ones.

If you mark all correct choices for each problem, it will give you all points of the problem; otherwise each wrong choice (marked) will result in "negative points" proportionally.

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Problem 1

1. (5p) Why is it so difficult to predict the timing behaviour of a computer system?
 - it may run software with too many lines of code
 - it may have shared resources
 - it is impossible to decide whether a program terminates or not
 - it may consume too much energy
 - it may have cache memory
 - it may not support direct memory access (DMA)
 - it may have a processor with branch prediction
2. (5p) The characteristics of RTOS are:
 - typical footprints: 10MB to 1GB
 - typical footprints: 0.2 to 1MB
 - deterministic and fast responsiveness
 - fast and time-sharing
 - possible to change scheduling policies
 - tasks are scheduled using EDF
3. (5p) Resource servers can be useful for
 - isolating applications of different types
 - hierarchically designing systems that may have many sub-systems
 - improving the average response times of soft real-time tasks
 - guaranteeing the deadlines of event-driven hard real-time tasks
4. (5p) When RMS is used, to ensure schedulability, the maximal resource utilization is bounded by:
 - 100%
 - 99%
 - $n * (2^{1/n} - 1)$ where n is the number of tasks
 - 69.3%
5. (5p) The advantages of static scheduling include:
 - it is flexible and easy to change
 - it can be used to implement predictable and deterministic systems
 - it can be used to implement complex systems with large task sets
 - it may use less memory than dynamic scheduling
 - it can be used to implement safety-critical systems
6. (5p) We know that the bandwidth is 1Mbps for a CAN bus whose length is 50 meters. What is the bandwidth for a CAN bus of 100 meters?
 - 3-4Mbps
 - 1-2Mbps
 - 0.2-0.5 Mbps
 - 0.1-0.15 Mbps
7. (5p) For a system with periodic tasks that are dependent on each other with input/output relations, the response times of tasks may vary. Is it possible to avoid this and thus avoid release jitters?
 - this is not possible at all

- yes, deliver the result of a job as soon as possible*
 - yes, deliver the result of a job as late as possible*
 - yes, deliver the result of a job at the end of each period*
8. (5p) Which statements are correct?
- EDF and DMS are non stable*
 - EDF is non stable; DMS is stable*
 - both DMS and EDF uses dynamic priorities*
 - EDF is optimal, DMS is not*
 - both EDF and DMS are optimal*
9. (5p) Assume that a set of periodic tasks with deadlines equal to periods is schedulable on a single processor system using Shortest Job First Scheduling Algorithm. Then the task set is also schedulable using
- RMS*
 - Non-Preemptive EDF*
 - EDF*
 - FIFO*
 - DMS*
10. (5p) When you use CAN to implement a distributed system, you should
- make sure that messages sent from different nodes have different priorities*
 - know that sending messages with the same identity by different nodes may lead to errors*
 - send messages with periodic tasks*
 - assign the priorities of messages according to the periods of the sending tasks*
 - make sure that the priority ordering of messages is given according to their id's with the lower id number, the higher priority for each message*
 - send messages without receiver's explicit address*

Problem 2 (10p) For a periodic task set with utilization 1, what would be the worst case response time of a task if EDF is used for scheduling?

- the period of the task*
- the worst case execution time of the task*
- any number larger than the worst case execution time of the task*
- it could be larger than the period of the task*
- any number less than the period of the task*
- any number between the worst case execution time and the period of the task*

Problem 3 (10p) For a set of sporadic sets (C_i, D_i, T_i) where deadline $D_i \leq T_i$. Identify the correct statements in the following:

- if the task set is feasible then $\sum C_i/T_i \leq 1$*
- if $\sum C_i/T_i \leq 1$, the task set is feasible*
- if the task set is feasible then $\sum C_i/D_i \leq 1$*
- if $\sum C_i/D_i \leq 1$, the task set is feasible*
- if $\sum C_i/D_i \leq 0.693$, the task set is schedulable with DMS*

- () if the task set is schedulable then $\sum C_i/D_i \leq 0.693$
- () if $\sum C_i/T_i \leq 0.693$, the task set is schedulable with DMS
- () if the task set is schedulable then $\sum C_i/T_i \leq 0.693$

Problem 4 (10p) Assume a set of 12 non-preemptive sporadic tasks whose computation times are all 0.27ms and whose periods are 1ms, 2ms, 2ms, 7ms, 8ms, 10ms, 18ms, 20ms, 30ms, 31ms, 43ms and 50ms respectively. Assume that RMS is used for scheduling, there are no jitters and no context-switch overheads. The worst-case response time for the task with period 10ms is bounded by:

- () 0.27ms, () 0.61ms, () 0.82ms, () 1.22ms, () 1.81ms, (X) 2.43ms, (X) 2.5ms, (X) 2.6ms
- (X) 2.7ms, (X) 2.8ms, (X) 3.1ms, (X) 3.51ms, (X) 4.32ms, (X) 5.67ms, (X) 7.02ms

Problem 5 (10p) What are the differences between BIP (Basic Priority Inheritance Protocol) and HLP (Immediate Priority Inheritance Protocol)?

- () BIP can avoid deadlocks, HLP can not
- (X) HLP can avoid deadlocks, BIP can not
- () BIP lets a task run at a higher priority whenever it locks a semaphore
- (X) HLP lets a task run at a higher priority whenever it locks a semaphore
- () HLP is easier to implement than BIP
- () Using HLP, the blocking time a higher priority task may suffer is the sum of the worst case execution times that tasks with lower priorities may run using the resources shared with the higher-priority task.
- (X) Using HLP, the blocking time a higher priority task may suffer is the max of the worst case execution times that tasks with lower priorities may run using the resources shared with the higher priority task.

Problem 6 (10p) Assume a set of event-driven tasks with computing times C_i and deadlines D_i (assume $C_i \leq D_i$). Both may arrive at any time but the minimal distance between two arrivals of each task is $T_i > D_i$. Design a polling server for the tasks with period T and capacity C such that no deadline is violated when RMS is used for scheduling. Assume $C \leq T$ and R is the response time estimated using RMS analysis when the server is considered as a periodic task. Identify correct statements:

- (X) $T \leq D_i$ is necessary
- () $T + T \leq D_i$ is necessary
- (X) $T + R \leq D_i$ is necessary and sufficient
- (X) $T + C \leq D_i$ is necessary
- (X) $\sum C_i \leq C$ is necessary
- (X) $\sum C_i \leq C$ and $T + T \leq D_i$ is sufficient
- () $T + T \leq D_i$ is sufficient

Problem 7 (10p) Assume 10 tasks each with execution time 2 and period 100. Assume a "large" task with execution time 100 and period 101. These 11 tasks will be scheduled on a multicore platform with 10 cores.

- the task set is schedulable with global RMS
- the task set is schedulable with global EDF
- the task set is schedulable with partitioned RMS
- the task set is schedulable with partitioned EDF
- (2)** what is the minimal number of cores needed for partitioned scheduling? (put your answer in brackets)

Problem 8 (10p) Assume a system with a set of DRT tasks described by digraphs whose workload are $dbf_i(t)$ (demand bound functions). Identify the correct ones out of the following statements:

- The system is feasible on a single-core platform if for all $t \geq 0$, $\sum dbf_i(t) \leq t$
- If the system is feasible on a single-core platform then for all $t \geq 0$, $\sum dbf_i(t) \leq t$
- The system is feasible on a two-core platform if for all $t \geq 0$, $\sum dbf_i(t) \leq 2t$
- If the system is feasible on a two-core platform then for all $t \geq 0$, $\sum dbf_i(t) \leq 2t$

Problem 9 (10p) Assume a set of periodic tasks with periods as deadlines and utilizations as follows: 0.95, 0.99, 0.5, 0.1, 0.6, 0.7, 0.5, 0.2, 0.4, 0.2, 0.4, 0.1

If you partition the tasks onto a multiprocessor platform (with job splitting if necessary) and then run the respective scheduling policy, what is the least number of processors you need?

- 8-9 with EDF
- 6-7 with EDF
- 4-5 with EDF
- 1-3 with EDF
- 9-10 with RMS
- 7-8 with RMS
- 5-6 with RMS
- 1-4 with RMS

Problem 10 (10p)

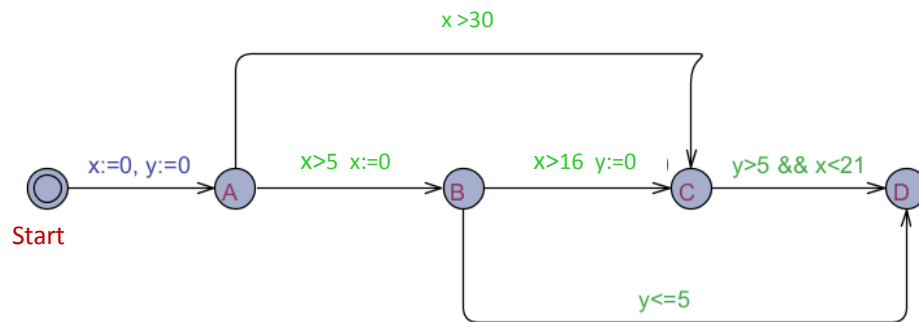


Figure 1. An Example of Timed Automaton where x and y are clocks and **Start** is the initial node.

Study the automaton in Figure 1. Identify the correct statements in the following:

- B is reachable and at B , it is always the case $y > 5$
- C is not reachable
- C is reachable and at C , it is always the case $x > 3$
- C is reachable and at C , it is possible $x > 100$
- C is reachable and at C , y can take any value larger than 0
- D is reachable
- D is reachable and at D , it is possible $x < 21$ and $y > 5$
- D is not reachable