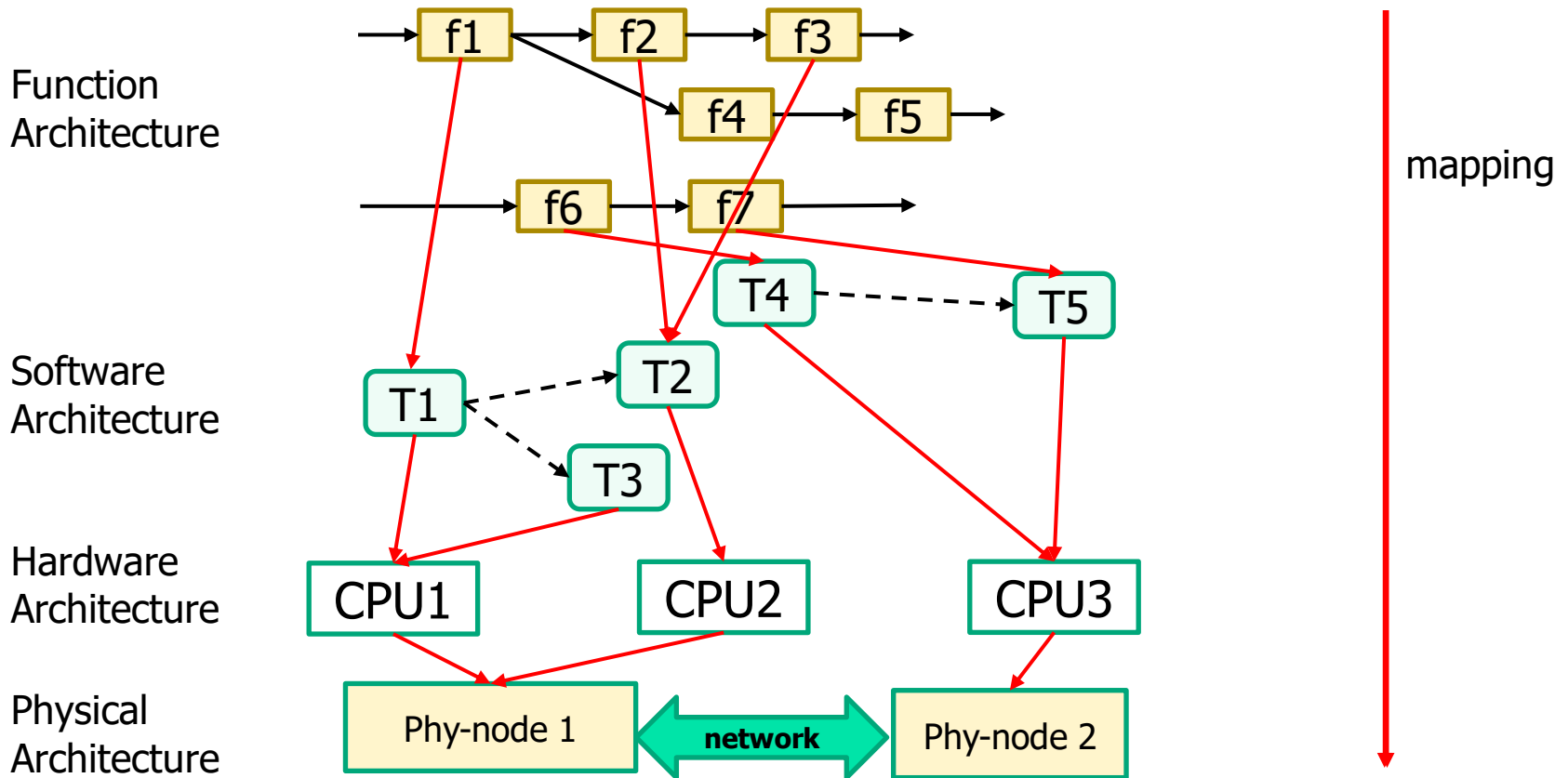


# Real Time Programming with Ada (1)

# Design of Real-Time Systems

(the **Autosar** approach in automotive industry)



# Real time programming

---

- It is mostly about **“Concurrent programming”**
- We also need to handle **Timing Constraints** on concurrent executions of tasks

However, remember:

- “concurrency” is a way of structuring computer programs e.g. “concurrent modules”: func 1, func 2 ...
- “concurrency” may be implemented on “uniprocessor platforms” using a scheduler: task 1, task 2 ...
- “concurrency” may be implemented on “multiprocessor platforms”: CPU 1, CPU 2 ...
- “concurrency” may be implemented on “distributed platforms”: Phy-Node 1, Phy-Node 2 ...

## This is the classic approach: **cyclic execution**

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- Program your tasks in any sequential language

```
loop
  run task 1
  run task 2
  run task 3
end loop
```

Efficient code, deterministic, predictable,

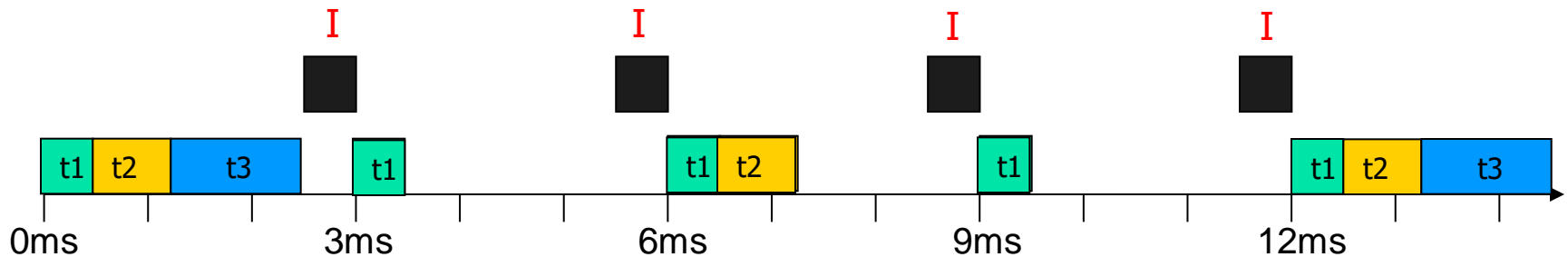
But (1) difficult to make it right, (2) difficult to reuse existing design

(3) extremely difficult for constructing large systems

# Cyclic Execution

Task	Period (rate)	Execution time
t1	3ms (333Hz)	0.5ms
t2	6ms (166Hz)	0.75ms
t3	12ms (83Hz)	1.25ms

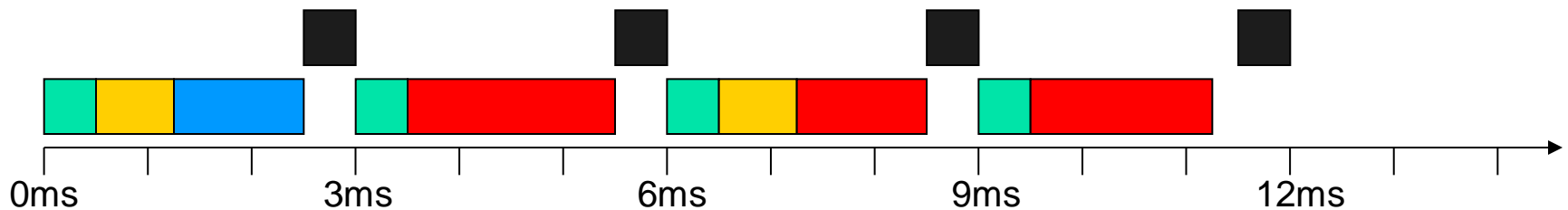
If there is an interrupt handling task: **I** with 0.5ms execution time every 3ms, we have to insert this ...



# Adding new functionality ... .. ?

---

- Every 12ms, 5.25ms is still free
- add t4 with period: 12ms & execution time: 5ms
- t4 has to be artificially partitioned



# Effect of new task at code level

```
void do_task_t4(void)
{
    /* Task functionality */
}
```



```
void do_task_t4_1(void)
{
    /* first bit */
    state_var_1 = x;
    state_var_2 = y;
    ...
}

void do_task_t4_2(void)
{
    x = state_var_1;
    ...
    /* second bit */
    state_var_3 = a;
    state_var_4 = b;
    ...
}

void do_task_t4_3(void)
{
    c = state_var_4;
    ...
    /* third bit */
}
```

```
int state_var_1;
int state_var_2;
int state_var_3;
int state_var_4;
```

```
void main(void)
{
    do_init();
    while(1) {
```

```
        do_task_t1();
        do_task_t2();
        do_task_t3();
        busy_wait();
        do_task_t1(); /* 3ms */
        do_task_t4_1();
        busy_wait();
        do_task_t1(); /* 6ms */
        do_task_t2();
        do_task_t4_2();
        busy_wait();
        do_task_t1(); /* 9ms */
        do_task_t4_3();
        busy_wait();
    }
```

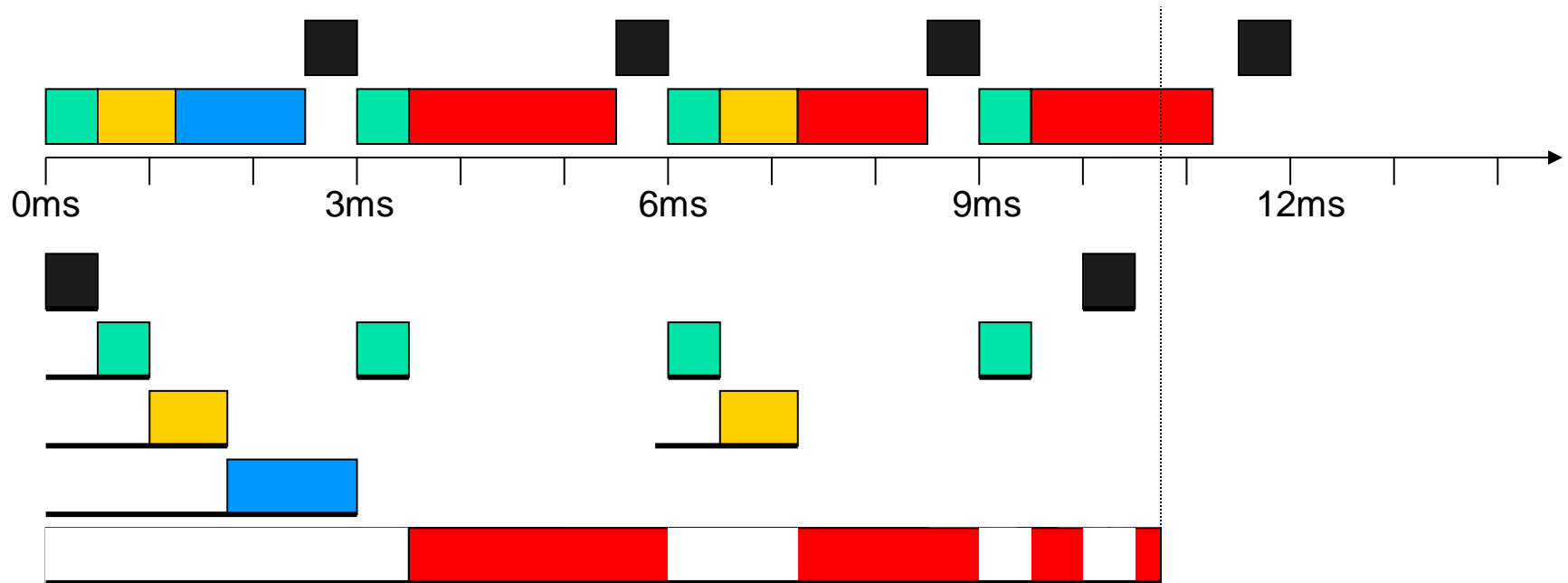
## This is “ad hoc”, but it is often used in industry

---

- This is not feasible for large software systems, say a few hundreds of control tasks
- This was why “Multitasking” came into the picture



# Cyclic Execution vs. Multitasking



# Ada95

---

- It is strongly typed OO language, looks like Pascal
- Originally designed by the US DoD as a language for large **safety critical systems** i.e. Military systems
  - Ada83
  - Ada95 + RT annex + Distributed Systems Annex
  - Ada 2005 (allows scheduling policies e.g. RR, EDF, dynamic priorities for protected types ...)

# The basic structures in Ada

---

- A large part in common with other languages
  - Procedures
  - Functions
  - Basic types: integers, characters, ...
  - Control statements: if, for, ..., case analysis
- **Any thing new?**
  - Abstract data type
    - Package (objects)
    - (protected package)
  - Concurrency
    - Tasking/Task declaration
  - Communication/synchronization
    - Rendezvous
    - Protected package/shared data type
  - Real Time facilities
    - Two pre-defined packages

# Typical structure of programs

---

Program Foo(...)

**Declaration 1** ←----- to introduce identities/variables  
and define data structures

**Declaration 2** ←----- to define "operations" : procedures, functions  
and/or tasks (concurrent operations)  
to manipulate the data structures

**Main program**

(Program body) ←----- a sequence of statements or "operations" to  
compute the result (output)

# Control-statements: If, case, for, while and assignment: $x := \text{exp}$

---

```
if TEMP < 15 then
    some smart code;
else
    do something else..;
end if;
```

```
case TAL is
    when <2 =>
        PUT_LINE("one or two");
    when >4 =>
        PUT_LINE("greater than 4");
end case;
```

```
for I in 1..12 loop
    PUT("in the loop");
end loop;
```

```
While B loop ... end loop
```

# Declarations and statements

---

- Before each block, you have to declare (define) the variables used, just like any sequential program

```
procedure PM (A : in INTEGER;  
             B: in out INTEGER;  
             C : out  INTEGER) is  
begin  
    B := B+A;  
    C := B + A;  
end PM;
```

## Types (like in Pascal or any other fancy languages)

---

type LINE\_NUMBER is range 1 .. 72

type WEEKDAY is (Monday, Tuesday, Wednesday);

type serie is array (1..10) of FLOAT;

type CAR is

record

REG\_NUMBER : STRING(1 .. 6);

MODEL : STRING(1 .. 20);

end record;

# Concurrent and Real-Time Programming with Ada

---

- **Abstract data types**
  - Package
  - (protected data package)
- **Concurrency: tasking**
  - task declaration
- **Communication/synchronization**
  - Rendezvous
  - protected package/shared data type
- **Real time facilities:**
  - Delay "time period" and Delay until "next-time point"
  - Real-time scheduling/"Fixed-priority scheduling"



---

Package --- Class/Object

# “Package”: abstract data type in Ada

---

- package definition ---- specification
- package body ---- implementation

# Package definition -- Specificaiton

---

- Objects declared in specification is visible externally.

```
package MY_PACKAGE is
-- declare/define data structures
  Type myobject is record
    Name: string
    Personalnr: integer
  End myobject
-- declare/define all public operations
  procedure myfirst_operation;
  procedure mysecond_operation;
  function mythird_operation (name: string) return myobject;
end MY_PACKAGE;
```

# Package body -- Implementation

- Implements package specification

(you probably want to use some other packages here e.g.. )

```
with TEXT_IO;
```

```
use TEXT_IO;
```

```
package body MY_PACKAGE is
```

```
    procedure myfirst_operation is
```

```
    begin
```

```
        myfirst_operation code here;
```

```
    end;
```

```
    function MAX (X,Y :INTEGER) return INTEGER is
```

```
    begin
```

```
        ... ..
```

```
    end;
```

```
    procedure mysecond_operation is
```

```
    begin
```

```
        PUT_LINE("Hello Im Ada Who are U");
```

```
        GET();
```

```
    end;
```

```
    function mythird_operation (name: string) return myobject is
```

```
    begin
```

```
        ... ..
```

```
    end;
```

```
end MY_PACKAGE;
```

# Protected data type

---

**protected** Buffer is

```
procedure read(x: out integer)
procedure write(x: in integer)
private
  v: integer := 0 /* initial value */
```

**protected body** Buffer is

```
procedure read(x: out integer) is
  begin x:=v end
procedure write(x: in integer) is
  begin v:= x end
```

(note that you can solve similar problems with semaphores)

# Concurrent and Real-Time Programming with Ada

---

- **Abstract data types**
  - package
  - (protected package)
- **Concurrency: tasking**
  - Task declaration
- **Communication/synchronization**
  - Rendezvous
  - Protected package/shared data type
- **Real time facilities:**
  - Delay "time period" and Delay until "next-time point"
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---

# Tasking

# Ada tasking: concurrent programming

---

- Ada provides at the language level light-weight tasks. These often referred to as threads in some other languages. The basic form is:

```
task T is                                ←----- specification
--- operations/entry (or simply: task T)
end T;
```

```
task body T is                            ←----- implementation/body
begin
---- processing----
end T;
```



## Example: the sequential case

---

```
procedure shopping is  
begin  
  buy-meat;  
  buy-salad;  
  buy-wine;  
end
```

Assume pre-defined procedures:  
 buy-meat  
 buy-salad  
 buy-wine

# The concurrent version

---

procedure shopping is

```
task get-salad;  
task body get-salad is  
begin  
buy-salad;  
end get-salad;
```

```
task get-wine;  
task body get-wine is  
begin  
buy-wine;  
end get-wine;
```

```
begin  
buy-meat;  
end
```

buy-salad and buy-wine  
will be activated concurrently  
here



And then run in parallel with  
buy-meat

# Creating Tasks

---

- Tasks may be declared at any program level
- Created implicitly upon entry to the scope of their declaration.
- Possible to declare task types to start several task instances of the same task type

# example

---

```
procedure Example1 is  
  task type A_Type;  
  task B;  
  A,C : A_Type;  
  
  task body A_Type is  
    --local declarations for task A and C  
  begin  
    --sequence of statements for task A and C  
  end A_Type;  
  
  task body B is  
    --local declarations for task B  
  begin  
    --sequence of statements for task B  
  end B;  
  
begin  
  --task A,C and B start their executions before the first statement of this procedure.  
end Example1;
```

# Task scheduling

---

- Allow priorities to be assigned to tasks in task definition
- Allow task dispatching policy to be set (Default: highest priority first)

task Controller is

```
    pragma Priority(10)
```

end Controller

## Task termination: a task will terminate if:

---

- It completes execution of its body
- It executes a terminate alternative of a select statement
- It is aborted:
  - `abort_statement ::= abort task_name {, task_name};`

---

# Communication/Synchronization

# Task communication/synchronization

---

- Message passing using "rendezvous"
  - **entry** and **accept**
- Shared variables
  - **protected objects/variables**



# Rendezvous

---

procedure foo

```
task T is  
  entry E(...in/out parameter...);  
end;
```

```
task body T is  
  begin  
    -----  
    accept E(... ..) do  
      ----- sequence of statements  
    end E;
```

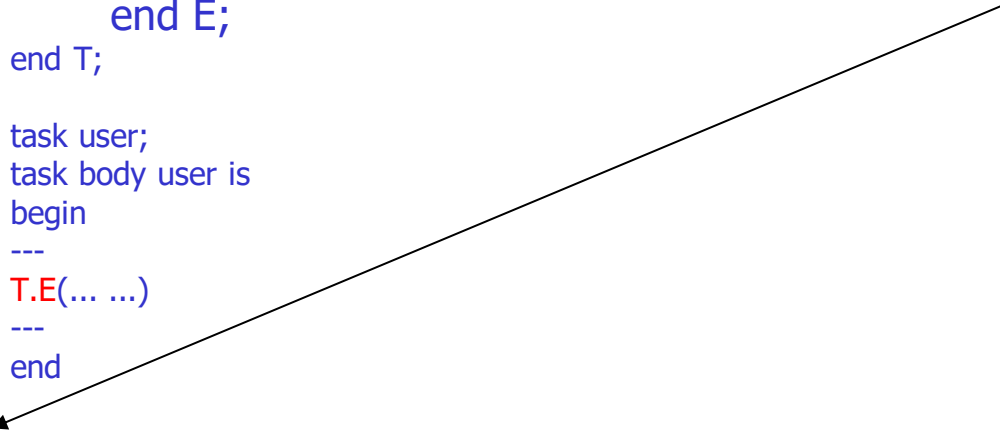
```
end T;
```

```
task user;  
task body user is  
begin  
---  
T.E(... ..)  
---  
end
```

```
begin
```

```
...  
end
```

T and user will be  
started concurrently



# Rendezvous

---

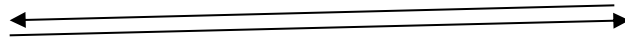
task body A is  
begin

...

B.Call;

...

end A



task body B is  
begin

...

accept Call do

....

end Call

...

end A

## This is implemented with Entry queues (the compiler takes care of this!)

---

- Each entry has a queue for tasks waiting to be accepted
  - a request to the entry is inserted in the queue
  - the first task in the queue will be "accepted" first (like the queue for a semaphore)
- By default, the queuing policy is FIFO
  - it can be different queuing policies

## Potential deadlocks

---

- Task A: .... B.b; accept a ...
- Task B: .... A.a; accept b ...

# An Example: Buffer

---

```
task buffer is
  entry put(X: in integer)
  entry get(x: out integer)
end;
```

```
task body buffer is
  v: integer;
begin
  loop accept put(x: in integer) do v:= x end put;
        accpet get(x: out integer) do x:= v end get;
  end loop;
end buffer;
```

```
---
buffer.put(...) ←----- other tasks (users)!!
Buffer.get(...)
----
```

# An Example, the Semaphore

---

- The Idea of a (binary) semaphore
- Two operations, p and v
  - p grabs semaphore or waits if not available
  - v releases the semaphore

# Program Semaphore using Task & RV. Synch.

---

## The specification

- task type Semaphore is
  - entry p;
  - entry v;
- end Semaphore;

## The implementation

- task body Semaphore is
  - begin
    - loop
      - accept p;
      - accept v;
    - end loop;
- end Semaphore;

Use a semaphore to protect critical sections

- Lock : Semaphore;
  - Lock.P;
  - Code for Critical Section
  - Lock.V;

# Choice: Select statement

---

```
task Server is
  entry S1(...);
  entry S2(...);
end Server;

task body Server is
  ...
begin
  loop
    --prepare for service
    select
      when <boolean expression> =>
        accept S1(...) do
          --code for this service
        end S1;
      or
        accept S2(...) do
          --code for this service
        end S2;
      or
        terminate;
    end select;
    --do any house keeping
  end loop;
end Server;
```



---

# Real-Time Facilities (next lecture)