



# Real Time Programming with Ada (1)

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## Real time programming

- It is mostly about "Concurrent programming"
- We also need to handle Timing Constraints on concurrent executions of tasks

However, remember:

- "concurrency" is a way of structuring computer programs  
e.g. three "concurrent modules": task 1, task 2 task 3
- "concurrency" is often implemented by "fast sequential computation" using a scheduler

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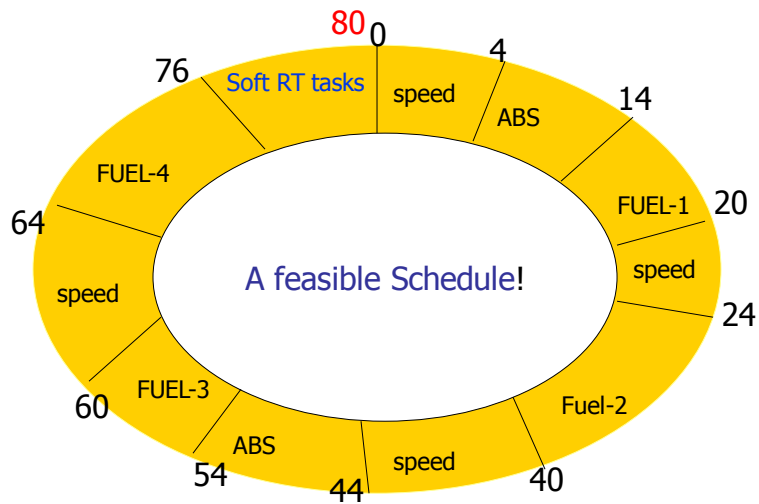
Programming the car controller

Process Speed: Loop next := get-time + 0.02 read sensor, compute, display... sleep until next End loop	Process ABS Loop next:=get-time + 0.04 Read sensor, compute, react sleep until next End loop
Process Fuel Loop next:=get-time + 0.08 read data, compute, inject ... sleep until next End loop	Soft RT Processes Loop read temperature elevator, stereo .... End loop

**Question:** do we need 4 CPUs to run these concurrently?

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Programming the car controller (3)



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## This is the classic approach: **cyclic execution**

- Program your tasks in any sequential language

```
loop
  do task 1
  do task 2
  do task 3
end loop
```

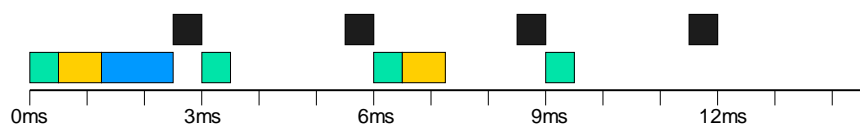
Efficient code, deterministic, predictable,  
But (1) difficult to make it right, (2) difficult to reuse existing design  
(3) extremely difficult for constructing large systems

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## Cyclic Execution

Task	Required sample rate	Processing time
t1	3ms (333Hz)	0.5ms
t2	6ms (166Hz)	0.75ms
t3	12ms (83Hz)	1.25ms

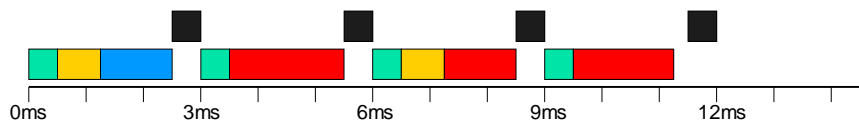
In addition, there is an interrupt handling task: **I**  
with 0.5ms processing time every 3ms



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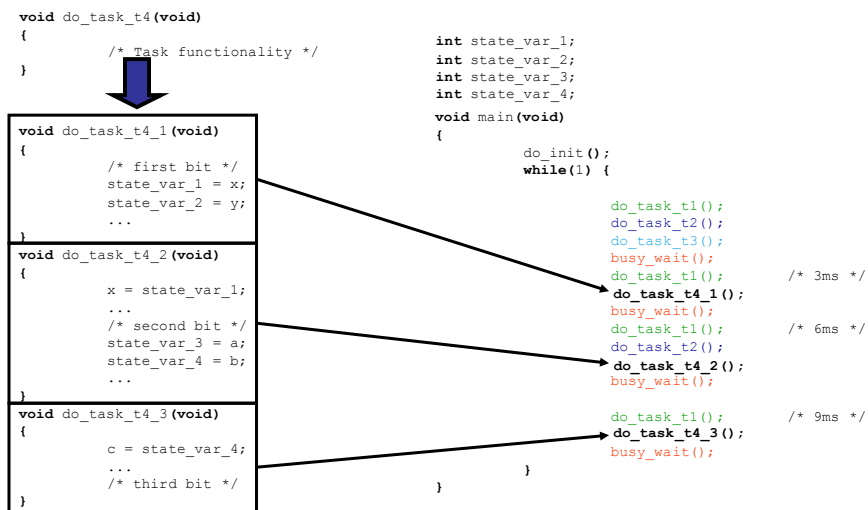
## Adding new functionality ... ..

- add task t4 with 12ms rate and 5ms processing time
- 12ms cycle has 5.25ms free time...
- ... t4 has to be artificially partitioned



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## Effect of new task at code level



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## This is "ad hoc", but it is often used in industry

- You just don't want to do this for large software systems, say a few hundreds of control tasks
- This was why "Multitasking" came into the picture

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## Concurrent programming with multitasking:

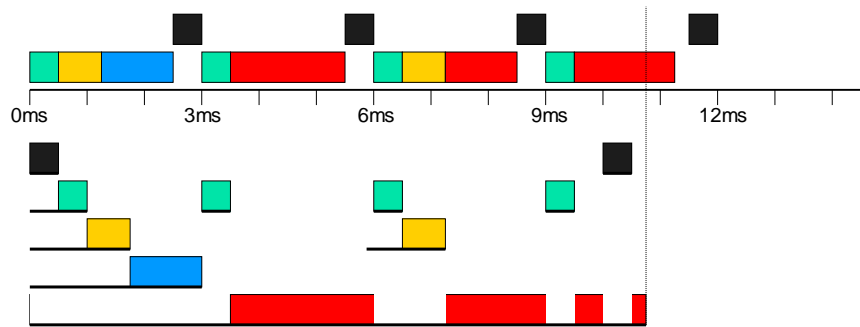
- Program your computation tasks, execute them concurrently with OS support e.g. in LegOS (or in Ada in slightly different syntax)

```
execi(foo1, ..., priority1, ...);  
execi(foo2, ..., priority2, ...);  
execi(foo3, ..., priority3, ...);
```

Will start three concurrent tasks running foo1, foo2, foo3

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## Cyclic Execution vs. Multitasking



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## Ada95

- It is strongly typed OO language, looks like Pascal
- Originally designed by the US DoD as a language for large **safety critical systems** i.e. Military systems
  - Ada83
  - Ada95 + RT annex + Distributed Systems Annex
  - Ada 2005 (allows scheduling policies e.g. RR, EDF, dynamic priorities for protected types ...)

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## The basic structures in Ada

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- A large part in common with other languages
  - Procedures
  - Functions
  - Basic types: integers, characters, ...
  - Control statements: if, for, ..., case analysis
- Any thing new?
  - Abstract data type: Packages (objects)
    - Protected data type
  - Tasking: concurrency
  - Task communication/synchronization: rendezvous
  - Real Time

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## Typical structure of programs

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Program Foo(...)

**Declaration 1** ←----- to introduce identities/variables  
and define data structures

**Declaration 2** ←----- to define "operations" : procedures, functions  
and/or tasks (concurrent operations)  
to manipulate the data structures

**Main program**  
(Program body) ←----- a sequence of statements or "operations" to  
compute the result (output)

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## Declarations and statements

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- Before each block, you have to declare (define) the variables used, just like any sequential program

```
procedure PM (A : in INTEGER;
             B: in out INTEGER;
             C : out  INTEGER) is
begin
  B := B+A;
  C := B + A;
end PM;
```

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## If, case, for: contrl-statements

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```
if TEMP < 15 then
  some smart code;
else
  do something else.;
end if;

case TAL is
  when <2 =>
    PUT_LINE("one or two");
  when >4 =>
    PUT_LINE("greater than 4");
end case;

for I in 1..12 loop
  PUT("in the loop");
end loop;
```

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## Types (like in Pascal or any other fancy languages)

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```
type LINE_NUMBER is range 1 .. 72
type WEEKDAY is (Monday, Tuesday, Wednesday);
type serie is array (1..10) of FLOAT;
```

```
type CAR is
  record
    REG_NUMBER   : STRING(1 .. 6);
    MODEL        : STRING(1 .. 20);
  end record;
```

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## Concurrent and Real-Time Programming with Ada

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- **Abstract data types**
  - package
  - protected data type
- **Concurrency**
  - Task creation
  - Task execution
- **Communication/synchronization**
  - Rendezvous
- **Real time:**
  - Delay "time period" and Delay until "next-time point"
  - Real-time scheduling/"Fixed-priority scheduling"

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## Package --- Class/Object

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## "Package": abstract data type in Ada

- package definition ---- specification
- Package body ---- implementation

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## Package definition -- Specificaiton

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- Objects declared in specification is visible externally.

```
package MY_PACKAGE is
-- declare/define data structures
  Type myobject is record
    Name: string
    Personalnr: integer
  End myobject
-- declare/define all public operations
  procedure myfirst_operation;
  procedure mysecond_operation;
  function mythird_operation (name: string) return myobject;
end MY_PACKAGE;
```

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## Package body -- Implementation

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- Implements package specification

```
(you probably want to use some other packages here e.g.. )
with TEXT_IO;
use TEXT_IO;

package body MY_PACKAGE is
  procedure myfirst_operation is
  begin
    myfirst_operation code here;
  end;
  function MAX (X,Y :INTEGER) return INTEGER is
  begin
    ... ..
  end;
  procedure mysecond_operation is
  begin
    PUT_LINE("Hello Im Ada Who are U");
    GET();
  end;
  function mythird_operation (name: string) return myobject is
  begin
    ... ..
  end;

end MY_PACKAGE;
```

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## Protected data type

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```
protected Buffer is
  procedure read(x: out integer)
  procedure write(x: in integer)
  private
    v: integer := 0 /* initial value */
```

```
protected body Buffer is
  procedure read(x: out integer) is
    begin x:=v end
  procedure write(x: in integer) is
    begin v:= x end
```

(note that you can solve similar problems with semaphores) 23

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## Tasking

## Ada tasking: concurrent programming

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- Ada provides at the language level light-weight tasks. These often referred to as threads in some other languages. The basic form is:

```
task T is                                ←----- specification
--- operations/entry (or simply: task T)
end T;
```

```
task body T is                            ←----- implementation/body
begin
---- processing----
end T;
```

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## Example: the sequential case

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```
procedure shopping is
begin
buy-meat;
buy-salad;
buy-wine;
end
```

Assume pre-defined procedures:  
buy-meat  
buy-salad  
buy-wine

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## The concurrent version

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procedure shopping is

```
task get-salad;  
task body get-salad is  
begin  
buy-salad;  
end get-salad;  
task get-wine;  
task body get-wine is  
begin  
buy-wine;  
end get-wine;  
begin  
buy-meat;  
end
```

buy-salad and buy-wine  
will be activated concurrently  
here

And then run in parallel with  
buy-meat

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## Creating Tasks

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- Tasks may be declared at any program level
- Created implicitly upon entry to the scope of their declaration.
- Possible to declare task types to start several task instances of the same task type

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## example

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```
procedure Example1 is
  task type A_Type;
  task B;
  A,C : A_Type;

  task body A_Type is
    --local declarations for task A and C
  begin
    --sequence of statements for task A and C
  end A_Type;

  task body B is
    --local declarations for task B
  begin
    --sequence of statements for task B
  end B;

begin
  --task A,C and B start their executions before the first statement of this procedure.
end Example1;
```

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## Task scheduling

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- Allow priorities to be assigned to tasks in task definition
- Allow task dispatching policy to be set (Default: highest priority first)

```
task Controller is
  pragma Priority(10)
end Controller
```

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## Task termination: a task will terminate if:

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- It completes execution of its body
- It executes a terminate alternative of a select statement
- It is aborted:
  - `abort_statement ::= abort task_name {, task_name};`

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## Communication/Synchronization

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## Task communication/synchronization

- Message passing using "rendezvous"
  - **entry** and **accept**
- Shared variables
  - **protected objects/variables**

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## Rendezvous

procedure foo

```
task T is
  entry E(...in/out parameter...);
end;

task body T is
  begin
    -----
    accept E(... ..) do
    ----- sequence of statements
    end E;
  end T;

task user;
task body user is
  begin
  ---
  T.E(... ..)
  ---
  end
end
begin
...
end
```

T and user will be started concurrently

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## Rendezvous

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```
task body A is
begin
...
B.Call;
...
end A

task body B is
begin
...
accept Call do
...
end Call
...
end B
```

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## This is implemented with Entry queues (the compiler takes care of this!)

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- Each entry has a queue for tasks waiting to be accepted
  - a call to the entry is inserted in the queue
  - the first task in the queue will be "accepted" first (like the queue for a semaphore)
- By default, the queuing policy is FIFO
  - it can be different queuing policies

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## An Example: Buffer

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```
task buffer is
entry put(X: in integer)
entry get(x: out integer)
end;

task body buffer is
  v: integer;
  begin
  loop accept put(x: in integer) do v:= x end put;
    accpet get(x: out integer) do x:= v end get;
  end loop;
end buffer;

---
buffer.put(...) ←----- other tasks (users)!!
Buffer.get(...)
----
```

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## Potential deadlocks

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- Task A: .... B.b; accept a ...
- Task B: .... A.a; accept b ...

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## An Example, the Semaphore

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- The Idea of a (binary) semaphore
- Two operations, p and v
  - p grabs semaphore or waits if not available
  - v releases the semaphore

## Program Semaphore using Task & RV. Synchron.

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### The specification

- task type Semaphore is
  - entry p;
  - entry v;
- end Semaphore;

### The implementation

- task body Semaphore is
  - begin
    - loop
      - accept p;
      - accept v;
    - end loop;
- end Semaphore;

## Program Semaphore using Task & RV

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- The implementation:
  - task body Semaphore is

```
begin
    loop
        accept p;
        accept v;
    end loop;
end Semaphore;
```

## Using the Semaphore

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- Declare an instance of a semaphore
  - Lock : Semaphore;
  - Now we can use Lock to protect critical sections

```
Lock.P;
    Code for Critical Section
Lock.V;
```

## Choice: Select statement

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```
task Server is
  entry S1(...);
  entry S2(...);
end Server;

task body Server is
  ...
begin
  loop
    --prepare for service
    select
      when <boolean expression> =>
        accept S1(...) do
          --code for this service
          end S1;
      or
        accept S2(...) do
          --code for this service
          end S2;
      or
        terminate;
    end select;
    --do any house keeping
  end loop;
end Server;
```

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## Real-Time Facilities (next lecture)

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