

Final Exam in Algorithms and Data Structures 1 (1DL210)

Department of Information Technology

Uppsala University

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Location: Polacksbacken

Time: 08:00 - 13:00

No books or calculator allowed.

Directions:

1. Do not write on the back of the paper
2. Do not use red ink
3. Write your **anonymous exam code** on each sheet of paper
4. **Important:** Unless explicitly stated otherwise, justify you answer carefully!
Answers without justification do not give any credits.

Good Luck!

Problem 1 (10p)

Order these functions in order of increasing asymptotic growth rate¹. If two of them have the same asymptotic growth rate, state that fact. No justification is needed.

$$\left(\frac{1000}{999}\right)^n \quad 2^n \quad \left(\frac{999}{1000}\right)^n \quad 2^{2 \cdot n} \quad n^2 \quad \log n \quad n^2 + n \log n \quad n$$

Problem 2 (10p)

For each of the propositions below, state whether that proposition is true or false. No justification is needed.

a) If $f(n) = \mathcal{O}(g(n))$ and $f(n) = \Omega(g(n))$, then we have $(f(n))^2 = \Theta((g(n))^2)$

b) If $f(n) = \mathcal{O}(g(n))$ and $f(n) = \Omega(g(n))$, then we have $f(n) = g(n)$

¹Here, log denotes the binary logarithm.

Problem 3 (15p)

The performance of QUICKSORT depends heavily on the choice of the so called *pivot* element. Assume that you have an implementation of QUICKSORT that simply chooses the leftmost element as pivot. Also assume that QUICKSORT returns the sorted array when finished.

- a) What is the asymptotic running time of QUICKSORT(A) in the *average case*?
- b) What is the asymptotic running time of QUICKSORT(QUICKSORT(A)) in the *average case*?

Problem 4 (10p)

- a) Draw the two input arrays for the *final* call to MERGE in MERGESORT([2, 4, 8, 3, 5, 1, 7, 6]).
- b) What is the worst case complexity of MERGE in MERGESORT? Note that the function takes as arguments two arrays A_1 and A_2 .
- c) What is the worst case complexity of MERGESORT?

Problem 5 (10p)

Assume that we have the set $S = \{14, 23, 32, 41, 50, 59, 68\}$ and we want to insert them into a hash table T of size at most 10, using chaining to resolve collisions.

- a) Provide a size of T and a suitable hash function h , such that
 - the distribution of elements in T by using h would be good for random input
 - h performs badly for the elements in S
- b) Provide a size of T and a suitable hash function h , such that

- the distribution of elements in T by using h would be good for random input
- h performs well for the elements in S

Justify each answer in at most 3 lines.

Problem 6 (15p)

- What is the largest possible height (i.e. number of levels) of a Binary Search Tree with n elements? Why?
- What is the worst case complexity of printing the elements of a Binary Search Tree in sorted order? Why?
- What is the *minimal* size of the rightmost subtree in a Max-Heap containing 10 elements?
- What is the *maximal* size of the rightmost subtree in a Max-Heap containing 10 elements?

Problem 7 (15p)

Give one possible BFS traversal (i.e. a sequence of nodes) starting from node A of the graph shown in Figure 1. Print the nodes *only* when they are finished. No justification needed.

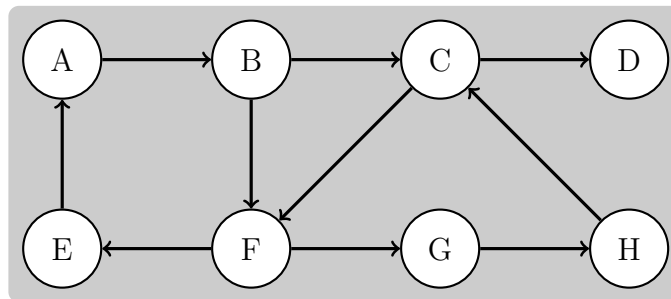


Figure 1: Graph for Problem 7

Problem 8 (15p)

Consider the algorithm VÄINÄMÖINEN, which operates on an array A containing the elements $A[0]$ to $A[n - 1]$:

```
VÄINÄMÖINEN( $A$ )
1   $n = \text{LENGTH}(A)$ 
2  for  $j = 0$  to  $n - 1$ 
3       $target = j$ 
4      for  $i = j + 1$  to  $n - 1$ 
5          if  $A[i] = A[target]$ 
6              return TRUE
7
8  return FALSE
```

- a) Describe, in 1 line, what VÄINÄMÖINEN does.
- b) Give a (tight) asymptotic upper bound on the average case running time of VÄINÄMÖINEN. (Is it $\mathcal{O}(1)$, $\mathcal{O}(\log n)$, $\mathcal{O}(n)$, etc?) Justify your answer in at most 5 lines.
- c) Propose a different way of doing the same thing that is asymptotically faster in the average case. Justify your answer in at most 5 lines.